Chapter 2

Noncommutative Gröbner Bases, and Projective Resolutions

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2.1 Overview

These notes consist of five sections. The aim of these notes is to provide a summary of the theory of noncommutative Gröbner bases and how to apply this theory in representation theory; most notably, in constructing projective resolutions.

Section 2.2 introduces both linear Gröbner bases and Gröbner bases for algebras. Section 2.3 surveys some of the basis algorithms of Gröbner basis theory. These include the Division Algorithm, and the Termination Theorem (Bergman's Diamond Lemma). Section 2.4 presents the noncommutative version of Buchberger's algorithm, universal Gröbner bases and considers the special case of finite dimensional algebras. Section 2.5 applies the theory of Gröbner bases to the study of modules. Projective presentations and resolutions are considered. A method of constructing projective resolutions for finite dimensional modules is given. Section 2.6 considers further theoretical applications. The study of Koszul algebras via Gröbner bases is presented.

2.2 Gröbner Bases

2.2.1 Linear Gröbner Bases In this section we consider only vector spaces. The ideas introduced here underlie much of what follows.

Throughout these notes, K will denote a field. Let V be a vector space. We fix a K-basis $\mathcal{B} = \{b_i\}_{i \in \mathcal{I}}$ where \mathcal{I} is an index set. One of the essential features of the theory of Gröbner basis is the selection of a well-ordered basis. Recall that > is a *well-order* on \mathcal{B} if > is a total order on \mathcal{B} and every nonempty subset of \mathcal{B} has a minimal element. The standard axioms of set theory imply that every set can be well-ordered.

Let > be a well-order on \mathcal{B} . We recall basic properties of >.

Proposition 2.1. [15] If \mathcal{B} is a set then > is a well-order on \mathcal{B} if and only if for each descending chain of elements of \mathcal{B} , $b_1 \ge b_2 \ge b_3 \ge \cdots$, there exists some N > 0, such that $b_N = b_{N+1} = b_{N+2} = \cdots$.

We will keep the following convention for the remainder of these lectures. If \mathcal{I} is an index set and $v_i \in V$ for $i \in \mathcal{I}$ then $\sum_{i \in \mathcal{I}} v_i$ implies that all but a finite number of $v_i = 0$. Thus, if I write $\sum_{i \in \mathcal{I}} \alpha_i b_i$ with $\alpha_i \in K$ and $b_i \in \mathcal{B}$, then all but a finite number of $\alpha_i = 0$.

One of the main uses of a well-order is to be able to have the notion of a largest basis element in a vector. If $v = \sum_{i \in \mathcal{I}} \alpha_i b_i$ we say b_i occurs in v if $\alpha_i \neq 0$. This leads to the following important definition.

Definition 2.1. If $\mathcal{B} = \{b_i\}_{i \in \mathcal{I}}$ is a basis of a vector space V and > is a well-order on \mathcal{B} , then if $v = \sum_{i \in \mathcal{I}} \alpha_i b_i$ is a nonzero element of V, we say b_i is the *tip of* v if b_i occurs in v and $b_i \ge b_j$ for all b_j occurring in v.

We denote the tip of v by Tip(v). If X is a subset of V we let

 $\operatorname{Tip}(X) = \{ b \in \mathcal{B} \mid b = \operatorname{Tip}(x) \text{ for some nonzero } x \in X \}.$

We let

$$\operatorname{NonTip}(X) = \mathcal{B} \setminus \operatorname{Tip}(X).$$

Thus, both $\operatorname{Tip}(X)$ and $\operatorname{NonTip}(X)$ are subsets of the fixed basis \mathcal{B} . Both sets are dependent of the choice of well order on \mathcal{B} .

It is not easy to see what the tip set of a subspace is from a generating set. Consider the following example where K is the field of rational numbers. Let $V = K^7$ and \mathcal{B} be standard basis ordered by $e_1 > e_2 > \cdots > e_7$. Let W be the subspace spanned by (1, 2, -1, 0, 2, 1, 5), (-1, -2, 0, 0, 1, -1, -3), (1, 2, -1, 0, 5, 1, 6). Then the tip set of W is $\{e_1, e_3, e_5\}$. Hence NonTip $(W) = \{e_2, e_4, e_6, e_7\}$.

We give a fundamental result which will be used often in what follows.

Theorem 2.1. Let V be a vector space over the field K with basis \mathcal{B} . Let > be a well-order on \mathcal{B} . Suppose that W is a subspace of V. Then

 $V = W \oplus Span(NonTip(W)).$

Proof. First we show that $W \cap \text{Span}(\text{NonTip}(W)) = (0)$.

Let $x \in \text{Span}(\text{NonTip}(W)) \setminus \{0\}$. If $x \in W$ then $\text{Tip}(x) \in \text{Tip}(W)$. But $\text{Tip}(x) \in \text{NonTip}(W)$ since $x \in \text{Span}(\text{NonTip}(W))$ and we would obtain a contradiction.

Now we show that W + Span(NonTip(W)) = V. This will use that > is a well-order on \mathcal{B} . Let $v \in V$ be such that Tip(v) is minimal with respect to the property that $v \notin W + \text{Span}(\text{NonTip}(W))$. We wish to show this leads to a contradiction. Consider Tip(v) = b. Let α be the coefficient of b in v. Note that

$$\operatorname{Tip}(v - \alpha \cdot b) < \operatorname{Tip}(v).$$

If $b \in \text{NonTip}(W)$ then by the above remarks and the minimality condition on $v, v - \alpha \cdot b = w + z$ with $w \in W$ and $z \in \text{Span}(\text{NonTip}(W))$. But since $b \in \text{NonTip}(W)$, we see

$$v = w + (z + \alpha \cdot b) \in W + \operatorname{Span}(\operatorname{NonTip}(W)),$$

a contradiction.

On the other hand, if $b \in \text{Tip}(W)$ let $w \in W$ such that Tip(w) = b. Let α be the coefficient of b in v and β be the coefficient of b in w. Then $v - (\alpha/\beta)w$ has smaller tip than v. Again, by the minimality condition on v,

$$v - (\alpha/\beta) \cdot w = w' + z,$$

where $w' \in W$ and $z \in \text{Span}(\text{NonTip}(W))$. Thus, we get a contradiction, since then

$$v = (w' + (\alpha/\beta)w) + z \in W + \text{Span}(\text{NonTip}(W)).$$

We now give the definition of a linear Gröbner basis. Let W be a subspace of V. We say a set of vectors $\mathcal{G} \subset W$ is a linear Gröbner basis for W with respect to > if

$$\operatorname{Span}(\operatorname{Tip}(\mathcal{G})) = \operatorname{Span}(\operatorname{Tip}(W)).$$

The reader should prove that if \mathcal{G} is linear Gröbner basis of W in V then $\text{Span}(\mathcal{G}) = W$.

Considering the fundamental Theorem 2.1, we see that every nonzero vector $v \in V$ can be written UNIQUELY as $w_v + N(v)$, where $w_v \in W$ and $N(v) \in$ Span(NonTip(W)).

Definition 2.2. We call N(v) the normal form of v.

Let the coefficient of the tip of a vector v be denoted $\operatorname{CTip}(v)$. There is a "best" linear Gröbner basis which can be defined as follows.

Definition 2.3. Let W be a subspace of V. We say a set \mathcal{G} of vectors in W is a reduced linear Gröbner basis for V (with respect to >) if the following conditions hold:

- 1. \mathcal{G} is a linear Gröbner basis of W.
- 2. If $g \in \mathcal{G}$ then $\operatorname{CTip}(g) = 1$.
- 3. If g and g' are distinct elements of \mathcal{G} then $\operatorname{Tip}(g) \neq \operatorname{Tip}(g')$.
- 4. If $g \in \mathcal{G}$ then $g \operatorname{Tip}(g) \in \operatorname{Span}(\operatorname{NonTip}(W))$.

The next result shows the existence and uniqueness of reduced linear Gröbner bases.

Proposition 2.2. Let V be a vector space and > a well-order on a basis \mathcal{B} of V. Let W be a subspace of V. Then there is a unique linear reduced Gröbner basis of V.

Proof. Let $\mathcal{T} = \text{Tip}(W)$. Then define $\mathcal{G} = \{t - N(t) \mid t \in \mathcal{T}\}$. By Theorem 2.1, $t - N(t) \in W$. Now, Tip(t - N(t)) = t since no basis element occuring in N(t) can be the tip of an element in W. Thus \mathcal{G} is a linear Gröbner basis of W. The rest of the properties of a reduced linear Gröbner basis are easy to check. \Box

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Note that by Proposition 2.3 below, a reduced linear Gröbner basis is in fact a basis. The next result gives both an algorithm to find the reduced linear Gröbner basis in the finite dimensional case and also, in conjunction with Proposition 2.2, proves the uniqueness of the reduced row echelon form of a matrix. We omit the details of row reduction techniques except to remind the reader that there are three row operations:

- 1. Multiply a row by a nonzero constant.
- 2. Interchange two rows.

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3. Change a row by adding a multiple of another row to it.

Given an $m \times n$ matrix M we view the m rows as vectors in K^n . The row space of M is the span of the row vectors. It is easy to see that performing row operations to a matrix does not change the row space.

Since we are fixing a K-basis of our vector space V, if $\dim_K V = n$ then we will identify V with K^n by using the ordering of \mathcal{B} . Let $\{e_1, \ldots, e_n\}$ be the standard basis for K^n and we assume that > is the order $e_1 > e_2 > \cdots > e_n$. Now if W is a subspace of V spanned by vectors w_1, \ldots, w_m then we view W as the row space of the matrix that has w_1, \ldots, w_m as the rows.

An $m \times n$ matrix M is reduced row-echelon form if

- 1. For some $1 \le r \le m$ the first r rows are nonzero and the last m r are zero.
- 2. There is an increasing sequence $1 \le c_1 < c_2 < \cdots < c_r \le n$ such the first nonzero column in the i^{th} row is c_i with entry 1.
- 3. If $i \neq j$ then entry in the j^{th} column of the i^{th} is 0.

The proof of the following proposition is left as an exercise.

Proposition 2.3. Let e_1, \ldots, e_n be the standard basis of K^n with well-order $e_1 > e_2 > \cdots > e_n$. If M is an $m \times n$ matrix is reduced row echelon form with r nonzero rows, then the reduced linear Gröbner basis for the row space M is given by the first r row vectors.

Note that as a consequence of the above discussion, using Theorem 2.1 and the uniqueness of the reduced Gröbner basis, we have an easy proof the uniqueness of the reduced row echelon form of a matrix.

2.2.2 Rings and Admissible Orders We now turn our attention to K-algebras. Let R be a K-algebra and let \mathcal{B} be a K-basis. We assume that \mathcal{B} is a semigroup with 0. That is, assume that under the multiplication of the ring, we have

 $b, b' \in \mathcal{B}$ implies $b \cdot b' \in \mathcal{B}$ or $b \cdot b' = 0$.

We call such a K-basis a multiplicative basis of R.

Examples:

- a): Let $R = K[x_1, ..., x_n]$, the commutative polynomial ring in *n* variables. Let $\mathcal{B} = \{\text{monomials}\}.$
- b): Let $R = K < x_1, \ldots, x_n >$ be the free associative algebra in *n* noncommuting variables. Let $\mathcal{B} = \{\text{monomials}\}.$
- c): Path algebras: Let $\Gamma = (\Gamma_0, \Gamma_1)$ be a finite directed graph. Here $\Gamma_0 = \{v_1, \ldots, v_n\}$ will be the vertex set (which we give some arbitrary order) and $\Gamma_1 = \{a_1, \ldots, a_m\}$ is the arrow set (which we also give some arbitrary order). Technically, we need two functions from $\Gamma_1 \to \Gamma_0$ corresponding to the origin vertex of the arrow and the terminus vertex of the arrow. We will denote these by o(a) and t(a) respectively. Furthermore, for vertices we set $o(v_i) = v_i = t(v_i)$.

Let \mathcal{B} denote the set of finite directed paths in Γ , including the vertices, which are viewed as paths of length 0. Each path p has a length, l(p), which is the number of arrows in p. We will write paths as follows:

$$p = a_{i_1}a_{i_2}\cdots a_{i_r}$$

where $a_{i_j} \in \Gamma_1$, $t(a_{i_j}) = o(a_{i_{j+1}})$ for $1 \le j \le r-1$. We let $o(p) = o(a_{i_1})$ and $t(p) = t(a_{i_r})$ and say that p is a (directed) path from $o(a_{i_1})$ to $t(a_{i_r})$. Of course in this case l(p) = r.

We give \mathcal{B} structure of a semigroup with 0 via concatenation. That is, if $p = a_1 \dots a_r \in \mathcal{B}$ and $q = b_1 \dots b_s \in \mathcal{B}$ where $a_i, b_j \in \Gamma_1$ then

$$p \cdot q = \begin{cases} a_1 \dots a_r b_1 \dots b_s & \text{if } t(a_r) = o(b_1) \\ 0 & \text{otherwise} \end{cases}.$$

Note that under the above definition, $v_i \cdot v_j$ is 0 if $i \neq j$ and is $v_i v_i = v_i$ if i = j.

The fact that \mathcal{B} is a semigroup with a relatively easy structure (for computational purposes) is one of the essential features of Gröbner bases. As we will see, reducing computations to \mathcal{B} is one of the underlying reasons that the techniques are so powerful.

We are now in a position to introduce the *path algebra* which we will denote by $K\Gamma$. As a K-vector space, $K\Gamma$ has K-basis \mathcal{B} . Thus the elements of $K\Gamma$ are K-linear combinations of paths and look like

$$x = \sum_{i=1}^r lpha_i p_i$$

where $\alpha_i \in K$ and p_i are paths in \mathcal{B} . We call the α_i the *coefficient of* p_i and say that p_i occurs in x if $\alpha_i \neq 0$. Having described the additive structure of $K\Gamma$ we need to give the multiplicative structure. For this, we define the multiplication on \mathcal{B} and then extend linearly to all of $K\Gamma$. If $p, q \in \mathcal{B}$ we define via the semigroup

multiplication in \mathcal{B} . Thus, $p \cdot q$ is 0 if the terminus of p is different than the origin of q, otherwise $p \cdot q$ is the concatenated path pq. We let K act centrally in $K\Gamma$ by $\alpha \cdot \sum \alpha_i p_i = \sum (\alpha \alpha_i) p_i$. The action of K commutes with all elements of $K\Gamma$. In particular, $K\Gamma$ is the semigroup algebra for the the semigroup \mathcal{B} .

Theorem 2.2. Basic Properties Let K be a field and Γ a finite directed graph. Then

1. $K\Gamma$ is a K-algebra with 1.

2.
$$1 = \sum_{v \in \Gamma_0} v$$

- 3. $\{v : v \in \Gamma_0\}$ is a full set of primitive orthogonal idempotents for $K\Gamma$. (A primitive idempotent is an idempotent that cannot be written as a sum of two orthogonal nonzero idempotents and we say two idempotents are orthogonal if their product, in either order, is 0)
- 4. $K\Gamma$ is a positively **Z**-graded ring with the elements of \mathcal{B} homogeneous and if $p \in \mathcal{B}$ has length l then $p \in (K\Gamma)_l$.
- 5. $K\Gamma$ is a tensor algebra.

Now assume that R is a K-algebra with multiplicative basis \mathcal{B} . We do not want an arbitrary well-order on \mathcal{B} . We want one that works well with the multiplicative structure of \mathcal{B} . For this, we introduce the following definition.

We say a well-order > on \mathcal{B} is *admissible* if it satisfies the following two conditions where $p, q, r, s \in \mathcal{B}$:

- 1. if p < q then pr < qr if both $pr \neq 0$ and $qr \neq 0$.
- 2. if p < q then sp < sq if both $sp \neq 0$ and $sq \neq 0$.
- 3. if p = qr then p > q and p > r.

Note that in case R is the commutative or noncommutative polynomial ring, the requirements that $pr \neq 0$, $qr \neq 0$, $sp \neq 0$, and $sq \neq 0$ all can be dropped. If R is a path algebra, then we can obtain 0 when multiplying basis elements.

The above properties restrict the orderings under consideration to those that "work well" with the multiplicative structure of \mathcal{B} . In the commutative theory of Gröbner bases, admissible orders are usually called "term orders" or "monomial orders".

We now give some examples which show that there are "natural" admissible orders on \mathcal{B} in the case of path algebras. For the following examples, assume we are given a finite directed graph Γ .

Example 2.1. The (left) length-lexicographic order:

Order the vertices and arrows arbitrarily and set the vertices smaller than the arrows. Thus

 $v_1 < \cdots < v_n < a_1 < \cdots < a_m.$

If p and q are paths of length at least 1, set p < q if l(p) < l(q) or if $p = b_1 \dots b_r$ and $q = b'_1 \dots b'_r$ with $b_i, b'_j \in \mathcal{B}$ and for some $1 \le i \le r, b_j = b'_j$ for j < i and $b_i < b'_i$.

The right length-lexicographic order is defined similarly.

Example 2.2. The (left) weight-lexicographic order:

Let $W: \Gamma_1 \to \{1, 2, 3, ...\}$ be a set map. Define $W: \mathcal{B} \to \{1, 2, 3, ...\}$ to be the natural extension. That is, W(v) = 0 if v is a vertex and if $a_i \in \Gamma_1$ define $W(a_1 \cdots a_r) = \sum_{i=1}^r W(a_i)$. Next, order the vertices and set the vertices smaller than the arrows. Order the arrows in such a fashion that if W(a) < W(b) then a < b.

Finally, define p < q if W(p) < W(q) or if W(p) = W(q) then use the left lexicographic order. Notice that the length-lexicographic order is a special case of the weight-lexicographic order. Mainly, give every arrow weight 1.

This order is sometimes called the *degree-lexicographic* order also.

It should be pointed out that the (left) lex order is NOT admissible in general. For example, if Γ has one vertex and two loops, a and b with b > a then we get

$$b > ab > aab > aaab > \cdots$$
.

Hence it is not a well-order. This is different from the commutative case.

There are some other, less obvious admissible orders.

Example 2.3. The (left) weight-reverse-lex order:

Take $W: \Gamma_1 \to \{1, 2, 3, ...\}$ and > define on the arrows and vertices as in the weight-lex order. Note that arrows and paths of positive length have positive weights.

Define p < q if W(p) < W(q) or if W(p) = W(q) then p > q in the right lex order. This is a well-order since there are only a finite number of paths of any weight.

One final example:

Example 2.4. The Total Lexicographic Order:

Label the arrows arbitrarily, say a_1, \ldots, a_m . Arbitrarily order the vertices and let them be smaller than any path of positive length. If p and q are paths, then p < q if there is some $i, 1 \le i \le m$, such that the number of a_j 's occuring in p and q are the same for j < i and the number of a_i 's occurring in p is less than the number of a_i 's occurring in q. If p and q have the same number of each arrow then p < q in lexicographic ordering (for some choice of ordering on the arrows).

2.2.3 Gröbner Bases Let R be a K-algebra with multiplicative basis \mathcal{B} and admissible order >. Let I be an ideal in R.

Definition 2.4. We say that a set $\mathcal{G} \subset I$ is a Gröbner basis for I with respect to > if

$$<$$
Tip $(\mathcal{G})>=<$ Tip $(I)>$.

That is, the two-sided ideal generated by the tips of \mathcal{G} equals the two-sided ideal generated by the tips of I. Equivalently, $\mathcal{G} \subset I$ is a Gröbner basis for I if for every $b \in \text{Tip}(I)$ there is some $g \in \mathcal{G}$ such that Tip(g) divides b; i.e., there are basis elements $p, q \in \mathcal{B}$ such that pTip(g)q = b.

We end this section by recalling Theorem 2.1 in the setting of rings. Suppose that R is K-algebra with multiplicative basis \mathcal{B} and admissible order >. Then if I is an ideal then

 $R = I \oplus \text{Span}(\text{NonTip}(I))$

as vector spaces. In particular, every nonzero element r of R can be written uniquely as $r = i_r + N(r)$ where $i_r \in R$ and $N(r) \in \text{Span}(\text{NonTip}(I))$. N(r) is still called the *normal form of* r. In the next section, I will address the existence of a reduced Gröbner basis and how to construct Gröbner bases.

2.3 Algorithms

2.3.1 Monomial Ideals and Reduced Gröbner Bases We begin with a proposition which shows that not all rings with a multiplicative basis admit an admissible ordering on the basis. Recall, if \mathcal{B} is a subset of R we say b_1 divides b_2 (in \mathcal{B}) if $b_1, b_2 \in \mathcal{B}$ and there are elements $c, d \in \mathcal{B}$ such that $b_2 = cb_1d$.

Proposition 2.4. If a multiplicative basis \mathcal{B} of a ring R admits an admissible order then every infinite sequence of elements of \mathcal{B} , b_1, b_2, b_3, \ldots , such that b_i divides b_{i-1} for $i \ge 1$ stabilizes; that is, for some N, $b_N = b_{N+1} = b_{N+2} = \cdots$.

Proof. The proof follows from the properties of an admissible order. Namely, suppose > is an admissible order for \mathcal{B} . Then, if b_i divides b_{i-1} , we get $b_1 \geq b_2 \geq b_3 \geq \cdots$. Since > is a well-ordering of \mathcal{B} , we get the desired result. \Box

Let R be a ring with multiplicative basis \mathcal{B} . We will call the elements of \mathcal{B} monomials. We say an ideal I in R is a monomial ideal if it can be generated by elements of \mathcal{B} . The next proposition is important in the definition of a reduced Gröbner basis.

Proposition 2.5. Let R be K-algebra with multiplicative basis \mathcal{B} which admits an admissible order. If I is a monomial ideal then I has a unique minimal monomial generating set. That is, there is a unique set of generators of I, none can be omitted and still generate I.

Proof. Consider the set A of all monomials in I. Let $M = \{p \in A \mid \text{ if } q \in A \text{ divides } p \text{ then } q = p\}$. Note that by Proposition 2.4, M is not empty. We claim that M is the unique minimal generating set of I. To show that M generates I, let B be a set of monomials that generates I. If $b \in B$, then for some $m \in M$, m divides b. Hence, $B \subset A$. Thus $I \subset A$.

If M' is another set of monomial set of generators of I, then every $m \in M$ is divisible by some $m' \in M'$. But then, by definition of M, m = m' and we have shown that $M \subset M'$.

Note that the minimal set of generators given in the above proposition is independent of any particular admissible order. The existence of an admissible order is necessary to show that M is nonempty. Furthermore, it is possible that the unique minimal monomial generating set is not finite. This differs from the commutative case in that Dickson's Lemma [15] proves that every monomial ideal in commutative polynomial ring is can be generated by a finite number of monomials. Good references for the commutative theory are [15, 9, 1].

Let I be an ideal in R and > an admissible order on a multiplicative basis \mathcal{B} . Then if \mathcal{G} is a Gröbner basis of I, then $\operatorname{Tip}(\mathcal{G})$ must contain the minimal monomial generating set of $< \operatorname{Tip}(I) >$.

If I is an ideal in R, we let I_{MON} be the ideal generated by Tip(I) (given some admissible order >). Then there is a unique minimal monomial generating set \mathcal{T} of I_{MON} .

Definition 2.5. The reduced Gröbner basis for I with respect to > is

$$\mathcal{G} = \{t - N(t) \mid t \in \mathcal{T}\}.$$

The proof of the next result is left as an exercise.

Proposition 2.6. Let > be an admissible order on a multiplicative basis \mathcal{B} of a K-algebra R. Let I be an ideal in R. Let \mathcal{G} be the reduced Gröbner basis for I. Then the following holds.

- 1. G is a Gröbner basis for I.
- 2. If $g \in \mathcal{G}$ then CTip(g) = 1.
- 3. If $g \in \mathcal{G}$ then $g Tip(g) \in Span(NonTip(I))$.
- 4. $Tip(\mathcal{G})$ is the minimal monomial generating set for I_{MON} .

2.3.2 The Division Algorithm In this section, we present a "division algorithm" in the rings we are studying. Throughout this section, we fix a K-algebra R with multiplicative basis \mathcal{B} and admissible order >.

Given an ORDERED set of elements $X = \{x_1, x_2, \ldots, x_n\}$ of R and another element $y \in R$ we show how to "divide" y by the set. We emphasize that the order of the elements affects the outcome of the division algorithm. What should division mean in this context? We mean that we find nonnegative integers m_1, \ldots, m_n and elements $u_{i,j}, v_{i,j}, r \in R$ for $1 \leq i \leq n$ and $1 \leq j \leq m_i$ such that

- 1. $y = \sum_{i=1}^{n} (\sum_{j=1}^{m_i} u_{i,j} x_i v_{i,j}) + r.$
- 2. $\operatorname{Tip}(y) \geq \operatorname{Tip}(u_{i,j}x_iv_{i,j})$ for all *i* and *j*.
- 3. For $b \in \mathcal{B}$ occurring in r, $\operatorname{Tip}(x_i)$ does not divide b for $1 \leq i \leq n$.

Note that it follows that $\operatorname{Tip}(r) \leq \operatorname{Tip}(y)$. We will call r a remainder by division by x_1, \ldots, x_n .

We now present the algorithm in pseudocode. Our presentation here takes the commutative division algorithm in [15] and changes it to make it noncommutative.

The Division Algorithm

```
INPUT:
               x_1,\ldots,x_n (ordered), y
OUTPUT:
                  m_1,\ldots,m_n,u_{i,j},v_{i,j},r
INITIALIZE:
                      m_1 := 0, \ldots, m_n := 0, r := 0, z := y, DIVOCCUR := False
WHILE
               (z \neq 0 \text{ and } \text{DIVOCCUR} == False)
                                                                    DO
    FOR
               (i = 1)
                            TO
                                    n
                                          DO
        IF
               \operatorname{Tip}(z) = u \operatorname{Tip}(x_i) v \text{ for } u, v \in \mathcal{B}
                                                               THEN
             m_i := m_i + 1
             u_{i,m_i} := [\operatorname{CTip}(z)/\operatorname{CTip}(x_i)]u
             v_{i.m_i} := v
             z := z - [\operatorname{CTip}(z)/\operatorname{CTip}(x_i)]ux_iv
              DIVOCCUR := True
        ELSE
                    i := i + 1
    IF
            DIVOCCUR == False
                                                THEN
         r := r + \operatorname{CTip}(z)\operatorname{Tip}(z)
         z := z - \operatorname{CTip}(z)\operatorname{Tip}(z)
        DONE
    DONE
DONE
```

We leave it to the reader to analyze the algorithm and show that it does what it is supposed to do. We give one small example.

Example 2.5. Let R be the noncommutative polynomial ring in three noncommuting variables x, y, z over a field K. Let \mathcal{B} be the set of monomials and > the length-lexicographic order with x > y > z. Let's divide $xy - x = f_1, xx - xz = f_2$ into zxxyx. Note that the tip of f_1 is xy and tip of f_2 is xx.

Beginning the algorithm, we see that $zxxyx = (zx)\text{Tip}(f_1)x$.

Thus $u_{1,1} = zx$, $v_{1,1} = x$ and we replace zxxyx by $zxxyx - zx(f_1)x = zxxx$. Now Tip (f_1) does not divide zxxx. Continuing, Tip (f_2) does. There are two ways to divide zxxx by xx and for the algorithm to be precise we must choose one. Say we choose the "left most" division. Then $zxxx = z(\text{Tip}(f_2))x$ and we let $u_{2,1} = z$, $v_{2,1} = x$ and replace zxxx by zxzx. Neither Tip (f_1) nor Tip (f_2) divide zxzx so we let r = zxzx and zxzx is replaced by 0 and the algorithm stops. We have

$$zxxyx = (zx)f_1x + zf_2x + zxzx.$$

The remainder is zxzx.

Note that if we interchange the order of f_1 and f_2 (so we start with f_2 first) we see the outcome of the division algorithm is

$$zxxyx = z(f_2)yx + zxzyx,$$

which gives a different remainder, namely, zxzyx.

Definition 2.6. If $X = \{x_1, \ldots, x_n\}$ (as an ordered set) and y is divided by X, we will denote the remainder r by $y \Rightarrow_X r$.

If one had an infinite set $X = \{x_1, x_2, ...\}$ with $\operatorname{Tip}(x_1) \leq \operatorname{Tip}(x_2) \leq \operatorname{Tip}(x_3) \leq \cdots$ and only a finite number of x_i 's with a given tip, we could perform division by X. That is, only a finite number of elements of X are needed for division and these can be determined prior to division. For, given y, there can only be a finite number of x_i 's such that $\operatorname{Tip}(x_i) \leq \operatorname{Tip}(y)$, say x_1, \ldots, x_n . By the desired properties of division, x_{n+1}, x_{n+2}, \ldots would never occur since we want $\operatorname{Tip}(y) \geq \operatorname{Tip}(u_{i,j}x_iv_{i,j})$. Thus, we may divide by an infinite set of elements of R provided that the set has the desired properties described earlier in the paragraph.

Now we show that when we have a Gröbner basis, the order of the x_i 's does not affect the remainder!

Proposition 2.7. If \mathcal{G} be a Gröbner basis for an ideal I in R. Let $y \in R$ and assume that $X = \{g_1, \ldots, g_n\} = \{g \in \mathcal{G} \mid Tip(g) \leq Tip(y)\}$. If $y \Rightarrow_X r$ then r is independent of the order of g_1, \ldots, g_n in X. In fact, r = N(y).

Proof. Consider $y \Rightarrow_X r$. Then, since $\operatorname{Tip}(r) \leq \operatorname{Tip}(y)$, we see that for each $g \in \mathcal{G}$, $\operatorname{Tip}(g)$ does not divide any basis element occurring in r. Hence $r \in \operatorname{Span}(\operatorname{NonTip}(I))$. Now $y = \sum_i \sum_j u_{i,j} g_i v_{i,j} + r$. But Theorem 2.1 implies that $y = i_y + N(y)$ with $i_y \in I$ and $N(y) \in \operatorname{Span}(\operatorname{NonTip}(I))$ unique. But $\sum_i \sum_j u_{i,j} g_i v_{i,j} \in I$ and $r \in \operatorname{Span}(\operatorname{NonTip}(I))$. Hence r = N(y).

Corollary 2.1. If \mathcal{G} is a Gröbner basis of an ideal I in R such that for each $b \in \mathcal{B}$ there are only a finite number of elements g in \mathcal{G} with $Tip(g) \leq b$ then there is an algorithm to find the normal form of elements of R.

In practice, Gröbner bases are found with only a finite number of terms with a given tip and hence the division algorithm gives an algorithm to find normal forms. Moreover, this is usually a "fast algorithm". The difficulty is in constructing Gröbner bases!

Note that we also now have a method to find the reduced Gröbner basis once we have found a Gröbner basis with only a finite number of terms with a given tip. This is an algorithm if I_{MON} has a finite set of monomial generators. The method proceeds as follows:

- 1. Given a Gröbner basis \mathcal{G} such that only a finite number terms of \mathcal{G} have a given tip.
- 2. Output will be the reduced Gröbner basis.

- 3. Find the minimal monomial generating set \mathcal{T} of I_{MON} .
- 4. For each $t \in \mathcal{T}$, using the division algorithm, calculate $t \Rightarrow_{\mathcal{G}} N(t)$.
- 5. The reduced Gröbner basis is $\{t N(t) | t \in \mathcal{T}\}$.

2.3.3 Overlap Relations and the Termination Theorem In this section, we give the noncommutative version of the S-polynomial found in commutative theory [15, 9]. We call these overlap relations. Throughout this section, R will be a K-algebra with multiplicative basis \mathcal{B} and admissible order > for \mathcal{B} .

Definition 2.7. Let $f, g \in R$ and suppose that there are elements $b, c \in \mathcal{B}$ such that

- 1. $\operatorname{Tip}(f)c = b\operatorname{Tip}(g)$.
- 2. Tip(f) does not divide b and Tip(g) does not divide c.

Then the overlap relation of f and g by b, c is

$$o(f, g, b, c) = (1/\operatorname{CTip}(f))fc - (1/\operatorname{CTip}(g))bg.$$

Note that $\operatorname{Tip}(o(f, g, b, c)) < \operatorname{Tip}(f)c = b\operatorname{Tip}(g)$. We give some examples to help clarify overlap relations.

Example 2.6. Let $R = K\Gamma$ where Γ is the graph:

$$\begin{array}{cccc} v_1 \bullet & \stackrel{a}{\longrightarrow} & \bullet v_2 \\ d \uparrow & & \downarrow b \\ v_4 \bullet & \stackrel{c}{\leftarrow} & \bullet v_3 \\ g \uparrow & & \downarrow e \\ v_6 \bullet & \stackrel{f}{\leftarrow} & \bullet v_5 \end{array}$$

Use the length-lexicographic order with $v_1 < \cdots < v_6 < a < b < c < \cdots < g$. Consider p = abcdabcd - abefgd and q = cdabcda - efgda. There are a number of overlap relations between p and q. We list them below.

- 1. $o(p,q,a,ab) = p \cdot a ab \cdot q = (-abefgd)a + abefgda = 0.$
- 2. $o(p,q,abcda,abcdab) = p \cdot abcda abcdab \cdot q$ = $-abefgdabcda + abcdabefgda \neq 0.$
- 3. $o(q, p, bcd, cd) = -efgdabcd + cdabefgd \neq 0$.
- 4. $o(q, p, bcdabcd, cdabcd) = -efgdabcdabcd + cdabcdabefgd \neq 0.$

As the above example shows, we need to look at all overlaps, including self overlaps; that is, overlaps of the form o(f, f, p, q). In this example, we have self overlaps o(p, p, abcd, abcd) and o(q, q, cdab, bcad). In the commutative case, one only uses the least common multiple of the leading monomials. In the noncommutative case, we must look at all overlaps.

Next we consider "tip reduction". Suppose that $X = \{x_1, \ldots, x_n\}$ is a set of elements in R. Let I be the ideal generated by X. If $\operatorname{Tip}(x_i) = u \operatorname{Tip}(x_j)v$ for some $u, v \in \mathcal{B}$, then letting

$$X' = \{x_1, \dots, x_{i-1}, x_i - [CTip(x_i)/CTip(x_j)]ux_jv, x_{i+1}, \dots, x_n\}$$

we that X' is a generating set of I also. Continuing in this fashion, we obtain a finite set of generators of I such that no tip of a generator divides the tip of another generator. This process is finite by the well-ordering assumption on >. Thus it is not unreasonable to assume that we have no tip divisions on a set of generators of an ideal.

Definition 2.8. We say a set of elements X is *tip reduced* if for distinct elements $x, y \in X$, Tip(x) does not divide Tip(y).

Before giving the termination Theorem, we need one more concept: uniformity.

Definition 2.9. We say an element $\sum_{i=1}^{n} \alpha_i b_i$, with $\alpha_i \in K^*$ and $b_i \in \mathcal{B}$, is *(left)* uniform if for each $c \in \mathcal{B}$, either $cb_i = 0$ for all $i, 1 \leq i \leq n$ or $cb_i \neq 0$ for all $i, 1 \leq i \leq n$.

Note that in the case of a noncommutative polynomial ring with $\mathcal{B} = \{\text{monomials}\}$, all elements are uniform. In the case of a path algebra, with $\mathcal{B} = \{\text{finite directed paths}\}$, then an element $\sum_{i=1}^{n} \alpha_i p_i$ is uniform if and only if there are vertices v, w such that for $1 \leq i \leq n$, the origin vertex of p_i is v and the terminus vertex is w. Note that every element $x \in K\Gamma$ is a sum of uniform elements. Since $1 = v_1 + \cdots + v_n$, we have $x = \sum_{i,j=1}^{n} v_i x v_j$ and we see that $v_i x v_j$ is uniform since all the paths that occur must have origin vertex v_i and terminus vertex v_j . It follows that every ideal in a path algebra can be generated by uniform elements.

We also note that if $f \in R$ is a uniform element and $b \in \mathcal{B}$ then both bf and fb are either 0 or uniform elements. We also remark that elements of \mathcal{B} are uniform.

We now give the termination Theorem which is a version of G. Bergman's Diamond Lemma [11, 16].

Theorem 2.3. Let R be a K-algebra with multiplicative basis \mathcal{B} and admissible order <. Suppose the \mathcal{G} is a set of uniform, tip reduced elements of R. Suppose for every overlap relation

$$o(g_1, g_2, p, q) \Rightarrow_{\mathcal{G}} 0,$$

with $g_1, g_2 \in \mathcal{G}$. Then \mathcal{G} is a Gröbner basis for $\langle \mathcal{G} \rangle$.

Proof. Assume that \mathcal{G} has the property that all overlap relations have remainder 0 under division by \mathcal{G} . Let $x \in I$ and we assume that tip(x) is not divisible by the tip of any element of \mathcal{G} . We show this leads to a contradiction. Without loss of generality, we amy assume that x is uniform. In this way, all multiplications

are in effect nonzero. Since we are assuming that \mathcal{G} is a generating set, we may write

$$x = \sum_{i,j} \alpha_{i,j} p_{i,j} g_i q_{i,j} \qquad (*)$$

where g_i varies over \mathcal{G} and $p_{i,j}, q_{i,j} \in B$. Consider all such ways of writing x. Let p^* be the largest path occuring on the rhs of (*). Since we are assuming that tip(x) is not divisible by the tip of any element of \mathcal{G} , by uniformity it follows that p^* is larger than tip(x) in the < order. Thus, the p^* 's terms cancel each other out.

Considering all ways of writing x as in (*), choose one such that p^* is as small as possible and has the fewest occurrences in the rhs of (*).

Since p^* does not occur on the lhs, it must appear in two summand of the rhs. So there exist i, j, i', j' so that

$$p^* = p_{i,j} tip(g_i)q_{i,j} = p_{i',j'} tip(g'_i)q_{i',j'}.$$

To simplify notation, write $p = p_{i,j}$, $g = g_i$, $q = q_{i,j}p' = p_{i',j'}$, $g' = g'_i$ and $q' = q_{i',j'}$. We proceed by a case by case study of the possible scenarios.

Case 1: length p < length p'.

In this case either, length $q \ge \text{length } q'$ or not.

Case 1.1: Length q < length q'.

Then tip(g') contains tip(g) and hence tip(g) divides tip(g') contradicting the hypothesis.

Case 1.2: length $q \ge$ length q'. We consider two possibilities.

Case 1.2.1: length $p' \ge \text{length } ptip(g)$.

Then there is no overlap of tip(g) and tip(g') in p^* . By the choice of lengths, it follows that there is a path q'' such that $p^* = ptip(g)q''g'q'$. Now if $g = \sum \alpha_i p_i + \alpha tip(g)$, and $g' = \sum \beta_i p'_i + \beta tip(p')$, then

$$pgq = pgq''(1/\beta)g'q' - pgq''(1/\beta)(g' - tip(g'))q'$$
$$= (\alpha/\beta)ptip(g)q''g'q' + \sum(\alpha_i/\beta)pp_iq''g'q' - \sum(\beta_i/\beta)pgq''(p'_i).$$

Thus, in writing pgq this way, can combine its tip with the tip of p'g'q' and lower the number of occurrances of p^* which contradicts the minimality assumption.

Case 1.2.2: length p' < length ptip(g).

Then there is an overlap of tip(g) and tip(g') is p^* . Say $tip(g)r = s \cdot tip(g')$. Thus $p^* = ptip(g)rq' = pstip(g')q'$. Then

$$pgq = c_g(tip(g))po(g,g',r,s)q + (c_g(tip(g))/c'_q(tip(g')))p'g'q'.$$

By assumption o(g, g', r, s) totally reduces over \mathcal{G} so is a K-linear combination of terms of the form $\hat{p}\hat{g}\hat{q}$ for paths $\hat{p}, \hat{q} \in B$ and $\hat{g} \in \mathcal{G}$, all of whose tips are smaller that $tip(g)r = s \cdot tip(g')$. So we may combine the occurrance of p^* in pgq with its occurrance in p'gq'. This again contradicts the minimality assumption.

Case 2: length p = length p'.

Then tip(g) divides tip(g') or vice versa - which contradicts the assumption on \mathcal{G} .

Case 3: length p >length p'.

Same as Case 1.

Note also that if division by \mathcal{G} of some overlap relation is NOT 0, then \mathcal{G} is not a Gröbner basis. This is easy to see since the remainder will have tip which is not divisible by the tip of any element of \mathcal{G} . Thus, when the hypothesis of Theorem 2.3 are met, we have an algorithm to determine if \mathcal{G} is a Gröbner basis (assuming \mathcal{G} is finite).

Example 2.7. Let $R = K\Gamma$ where Γ is the graph:

	a		ь	
$\overset{v_1}{\bullet}$	\rightarrow	$\overset{v_2}{\bullet}$	\rightarrow	v_3 \bullet
	←		\leftarrow	
	d		с	

Let $>_1$ be the length-lexicographic order with $v_1 < v_2 < v_3 < d < c < b < a$. Consider $\mathcal{G} = \{cdab - cb, bc - da\}$. Both elements are uniform and tip reduced. Then

$$o(bc - da, cdab - cb, dab, b) = -dadab + bcb \Rightarrow_{\mathcal{G}} - dadab + dab \neq 0.$$

Hence \mathcal{G} is not a Gröbner basis with respect to $>_1$.

But if we change $>_1$ to $>_2$ which is the length-lexicographic order with $v_1 < v_2 < v_3 < a < b < c < d$ then Tip(cdab-cb) = cdab and Tip(bc-da) = da. Thus the set is not tip reduced. If we tip reduce \mathcal{G} we get $\mathcal{G}' = \{cbcb-cb, -da+bc\}$. Then Tip(cbcb-cb) = cbcb and Tip(-da+bc) = da. The only overlap relation is

$$o(cbcb - cb, cbcb - cb, cb, cb) = 0.$$

Hence all overlap relations have remainder 0 and we conclude that under $>_2$, \mathcal{G}' is a Gröbner basis.

2.4 Computational Aspects

2.4.1 Construction of Gröbner bases In this section we present the noncommutative analog of Buchberger's algorithm [13] for constructing Gröbner bases. Given uniform elements f_1, \ldots, f_n in R, let $I = \langle f_1, \ldots, f_n \rangle$. Hence, by construction, I is generated by a set of uniform elements. The algorithm produces a (possibly infinite) sequence of uniform elements g_1, g_2, \ldots where $g_i = f_i$ for $1 \leq i \leq n$ and, for i > n, $g_i \in I$ such that

$$\operatorname{Tip}(g_i) \not\in < \operatorname{Tip}(g_1), \operatorname{Tip}(g_2), \ldots, \operatorname{Tip}(g_{i-1}) > .$$

It can be shown that $\{g_1, g_2, \ldots, g_m, g_{m+1}, \ldots\}$ is in fact a Gröbner basis for I. We present the algorithm in pseudocode.

```
f_1,\ldots,f_n
INPUT:
OUPUT:
                  q_1, q_2, q_3, \cdots
          i = 1 \text{ TO } n \text{ DO}
FOR
     g_i := f_i
     \mathcal{G} := \{g_1, \ldots, g_n\}
     Count := n
DO
     \mathcal{H} := \mathcal{G}
     FOR each pair of elements h, k \in \mathcal{H} AND each overlap relation of
     h, k
          DO
          IF
                  o(h, k, p, q) \Rightarrow_{\mathcal{H}} r \text{ AND } r \neq 0
                                                                     DO
                Count := Count + 1
               g_{\text{Count}} = r\mathcal{G} := \mathcal{G} \cup \{g_{\text{Count}}\}
          DONE
     DONE
WHILE
                 (\mathcal{H} \neq \mathcal{G})
```

Modifying the proof of the termination Theorem of the previous section, it can be shown that if $b \in \mathcal{B}$ is a minimal monomial generator of I_{MON} then for some m, $\text{Tip}(g_m) = b$. From this we get the next result.

Proposition 2.8. If I_{MON} has a finite set of monomial generators then the above algorithm terminates in a finite number of steps and yields a finite Gröbner basis.

Proof. If I_{MON} has a finite set of mononial generators, then the unique minimal monomial generating set must be finite. Suppose $\mathcal{T} = \{t_1, \ldots, t_s\}$ is the finite generating set of minimal monomials for I_{MON} . Then, by the remarks

preceeding the proposition, $\mathcal{T} \subset {\text{Tip}(g_1), \ldots, \text{Tip}(g_N)}$ for sufficiently large N. But then $\text{Tip}({g_1, \ldots, g_N})$ generates I_{MON} and hence is a Gröbner basis of I. But then division by ${g_1, \ldots, g_N}$ can only have remainder 0 since all overlap relations are elements of I. Thus the algorithm terminates in a finite number of steps outputting a Gröbner basis. \Box

For a discussion of improvements of the Buchberger algorithm in the commutative case we refer to [9]. Most of the discussion there can be translated to the noncommutative case.

2.4.2 Basic Computational Use of Gröbner Bases We are interested in studying quotient rings R/I. Elements of R/I consist of equivalence classes of the form f + I where $f \sim g$ if and only if $f - g \in I$. The K-algebra has addition and multiplication given by

$$(f + I) + (g + I) = ((f + g) + I)$$

 $(f + I) \cdot (g + I) = ((fg) + I)$

To study R/I we need a way studying the equivalence classes f+I. Assuming we have R represented on a computer via the given multiplicative basis \mathcal{B} , we would like to be able to find "good" representatives of equivalence classes. But given an admissible order > on \mathcal{B} we have such representatives.

Proposition 2.9. Let R be a K-algebra with multiplicative basis \mathcal{B} and admissible order > on \mathcal{B} . Let I be an ideal in K.

- 1. f + I = g + I if and only N(f) = N(g).
- 2. f + I = N(f) + I.
- 3. The map $\sigma : R/I \to R$ with $\sigma(f+I) = N(f)$, is a vector space splitting to the canonical surjection $\pi : R \to R/I$.
- 4. σ is a K-linear isomorphism between R/I and Span(NonTip(I)).
- 5. Identifying R/I with Span(NonTip(I)), then NonTip(I) is a K-basis of R/I contained in \mathcal{B} .

Proof. The first two parts are immediate consequences of Theorem 2.1 which states that $R = I \oplus \text{Span}(\text{NonTip}(I))$. The remaining parts are left as an exercise.

Thus, normal forms solves the problem of finding representatives of the equivalence classes. Furthermore, addition of two equivalence classes (f + I) + (g + I) is given simply by N(f) + N(g) = N(f + g) since Span(NonTip(I)) is a linear subspace of R. By Proposition 2.9, since $N(f) \in f + I$ we see that multiplication of classes is given by

$$N(N(f) \cdot N(g)),$$

that is, to find the representative of $(f+I) \cdot (g+I)$ simply multiply N(f)N(g) in R and then take the normal form of the result.

As pointed out in the last section, once a Gröbner basis has been found, the division algorithm provides an algorithm to find normal forms. In this way, one can use the computer to study quotient rings R/I.

2.4.3 Finite Dimensional Algebras Suppose that R/I is finite dimensional. What further can be said in this case?

Proposition 2.10. Let R be a finitely generated K-algebra with multiplicative basis \mathcal{B} and admissible order >. Suppose that I is an ideal such that $\dim_K(R/I) = N$. Then I_{MON} has a finite set of monomial generators.

Proof. Since R/I is isomorphic to Span(NonTip(I)) as vector spaces, it follows that NonTip(I) is a finite set since it is a basis of Span(NonTip(I)). Next, since R is finitely generated as an algebra, \mathcal{B} is finitely generated as a semigroup.

Let $X = \{b_1, \ldots, b_k\}$ generate \mathcal{B} . We show

 $\{bc \mid b \in X \text{ and } c \in \operatorname{NonTip}(I)\} \cap \operatorname{Tip}(I)$

generates I_{MON} . Suppose that t is a element of the minimal monomial generating set of I_{MON} . Then $t = b_{i_1}b_{i_2}\cdots b_{i_l}$ with $b_{i_j} \in X$. By minimality, $b_{i_2}\cdots b_{i_l} \notin$ Tip(I). Thus, $b_{i_2}\cdots b_{i_l} \in \text{NonTip}(I)$. Hence $t = b_{i_1}c$ with $c \in \text{NonTip}(I)$. This completes the proof.

The above result has the following immediate consequence.

Corollary 2.2. Let R be a finitely generated K-algebra with multiplicative basis and admissible order >. Suppose that I is an ideal generated by uniform elements and R/I is finite dimensional. Then I has a finite uniform Gröbner basis with respect to > and can be computed in a finite number of steps by the above algorithm.

Open Question: Given R and a multiplicative basis \mathcal{B} . Find necessary and sufficient conditions on an ideal I in R such that there is some admissible order > for which I has a finite Gröbner basis.

2.4.4 Universal Gröbner Bases If R is a K-algebra with multiplicative basis \mathcal{B} . Assume that R is finitely generated and hence so is B. Fix a set of generators of the semigroup $\mathcal{B}, b_1, \ldots, b_n$.

Definition 2.10. If $b \in \mathcal{B}$ we define the *length of b* to be k if b is a product of k generators and cannot be written as a product of fewer than k generators. We denote the length of b by length(B).

Thus length $(b_i) = 1$ for i = 1, ..., n and if $b \in \mathcal{B}$, length $(b) \ge 1$. Let I be an ideal in R.

Definition 2.11. We say a set of elements \mathcal{G} in I is a *universal Gröbner basis for* I if for every admissible order >, \mathcal{G} is a Gröbner basis with respect to >.

In some cases, finite universal Gröbner bases exist.

Example 2.8. Let R be the noncommutative polynomial ring in three noncommuting variables x, y, z over the rational numbers. Let \mathcal{B} be the set of monomials. Let I be the ideal generated by xy - 2yx, zx, zy. This set is a universal Gröbner basis since all overlap relations reduce to 0 whether the tip of xy - 2yx is xy or is yx. Hence the order doesn't affect that the set is a Gröbner basis. Note that R/I is not even noetherian.

In general, it is not known which ideals have finite universal Gröbner bases. One of the difficulties with this problem in the noncommutative case is that admissible orders are not classified.

We will show that finite universal Gröbner bases exist if R/I is finite dimensional. For this we introduce some terminology. If V is a vector space with basis \mathcal{B} , then the support of a vector v, $\operatorname{supp}_{\mathcal{B}}(v)$, is defined to be

 $\operatorname{supp}_{\mathcal{B}}(v) = \{ b \in \mathcal{B} \, | \, b \text{ occurs in } v \}.$

If X is a subset of V, the support of X is

$$\operatorname{supp}_{\mathcal{B}}(X) = \{\operatorname{supp}_{\mathcal{B}}(x) \, | \, x \in X\}.$$

Proposition 2.11. Let R be a finitely generated K-algebra with multiplicative basis \mathcal{B} which admits an admissible order. Let I be an ideal such that R/I is finite dimensional. Then there exists a finite universal Gröbner basis for I.

Proof. Let $d = \dim_K(R/I)$ and let $\pi: K\Gamma \to K\Gamma/I$ be the canonical surjection. Assume that \mathcal{B} is generated by $X = \{b_1, \ldots, b_n\}$. We claim that every $b \in \mathcal{B}$ with of length longer than D = (d+2)n has a factor which is the tip of an element in I of smaller length. That is, if the length of b is greater than D, then there is some monomial q of smaller length such that b = sqt and such that q is the tip of some element in I. Suppose b is a monomial of length longer than D. Write $b = b_{i_1}b_{i_2}\ldots b_{i_E}$ with $b_{i_j} \in X$ and E > D. Then some generator of \mathcal{B} , say b^* , must occur at least d + 2 times. Hence we have a factorization of b into monomials, $b = p_1c_1\cdots c_{d+1}p_2$, where each c_i is of the form $b^*b_{i_j}\ldots b^*$ and is of length at least 1 but of length less than the length of b. Since $\dim_K(R/I) = d$, the set $\{\pi(c_1), \ldots, \pi(c_{d+1})\}$ is linearly dependent over K. Therefore there is a nontrivial linear combination $x = \sum_{j=1}^{d+1} \alpha_i c_i \in I$. Thus, some c_i is a tip in the given ordering and we have proven the claim.

Now let V be the vector subspace of R with basis Ω consisting of monomials of length bounded by D. Let Y be the subspace $I \cap V$ of V. Since V is finite dimensional, the support of Y is a finite set, say $\operatorname{supp}_{\Omega}(Y) = \{s_1, \ldots, s_m\}$. For each $s_i \in \operatorname{supp}_{\Omega}(X)$, choose an element $f_i \in I$ such that $\operatorname{supp}(f_i) = s_i$. We now show that $\mathcal{G} = \{f_1, \ldots, f_m\}$ is a universal Gröbner basis for I.

Let > be an admissible order and let $h \in I$. Suppose the remainder h^* of division of h by \mathcal{G} is not 0. No monomial occurring in h^* is divisible by any $\operatorname{Tip}(f_i), i = 1, \ldots, m$. Let h^* have that property that the longest monomial in h^* is minimal in the set of all $h \in I$ such that no monomial in h is divisible by any $\operatorname{Tip}(f_i)$. Let C be the length of the longest monomial in h^* . If $C \leq D$

then $\operatorname{supp}(h^*) \in \operatorname{supp}_{\Omega}(Y)$ and hence there is some $f_i \in \mathcal{G}$ such that $\operatorname{supp}(f_i) = \operatorname{supp}(h^*)$. But then $\operatorname{Tip}(f_i) = \operatorname{Tip}(h^*)$, a contradiction. If C > D, then by an earlier argument, if $p \in \mathcal{B}$ has length greater than D and occurs in h^* , then there is some $f \in \mathcal{G}$ such that $\operatorname{Tip}(f)$ is a subpath of p. Thus, p can be reduced by an element in \mathcal{G} . But this contradicts our assumption on h^* and we are done. \Box

2.5 Modules, Presentations and Resolutions

In this section I will address the problem of studying R/I-modules. If R/I is not finite dimensional and the module is not finite dimensional, then the only possibility of computationally handling this situation is using generators and relations; a setup ideally suited for the use of the theory of Gröbner bases. Throughout this section we fix a K-algebra R, a multiplicative basis \mathcal{B} , and an admissible order >. Also fix an ideal I in R and assume we have a Gröbner basis \mathcal{G} of I with respect to >.

2.5.1 Modules Suppose that M is a *finitely presented* right R/I-module. This means that M is the cokernel of a map between finitely generated projective R/I-modules. This is equivalent to assuming that M is the cokernel of a map between two finitely generated free R/I-modules.

Let $f_1: F_1 \to F_0$ be an R/I-homomorphism between finitely generated free right R/I-modules. We choose bases for F_0 and F_1 , say $\{e_1, \ldots, e_n\}$ for F_0 and $\{d_1, \ldots, d_m\}$ for F_1 . Then f_1 is given by an $m \times n$ matrix $A_{f_1} = (a_{i,j})$ where $a_{i,j} \in R/I$. Then

$$f_1(d_i) = \sum_j^n e_j a_{i,j}.$$

Let $M = \operatorname{Coker}(f_1) = F_0 / \operatorname{Im}(f_1)$ and let $f_0 : F_0 \to M$ be the canonical map.

Note that computationally, we can represent the matrix A_{f_1} by an $m \times n$ matrix with entries in Span(NonTip(I)).

Before specializing, I present a way of viewing M as an ideal in a ring where we can use Gröbner basis theory to study M. In the representation theory of finite dimensional algebras, this is called "one point extension".

Let $\bar{S} = \begin{pmatrix} K & M \\ 0 & R/I \end{pmatrix}$. Note that multiplication is just matrix multiplication; namely,

$$\left(\begin{array}{cc}k&m\\0&r\end{array}\right)\cdot\left(\begin{array}{cc}k'&m'\\0&r'\end{array}\right)=\left(\begin{array}{cc}kk'&km'+mr'\\0&rr'\end{array}\right),$$

which makes sense in that we view M as a K-R/I-bimodule (with elements of K commuting will elements of M).

We now show how to view \overline{S} as quotient of a K-algebra for which we can find a multiplicative basis and an admissible order related to that of R.

Let $G_0 = \prod_{i=1}^n R$ with basis e_1^*, \ldots, e_n^* .

Define S to be the K-algebra by $S = \begin{pmatrix} K & G_0 \\ 0 & R \end{pmatrix}$. Set $\pi : S \to \overline{S}$ by

$$\pi\left(\begin{pmatrix} k \sum_{i=1} e_i^* r_i \\ 0 & r \end{pmatrix}\right) = \begin{pmatrix} k f_0(\sum_{i=1} e_i^* N(r_i)) \\ 0 & N(r) \end{pmatrix}.$$

Note that we have identified R/I with Span(NonTip(I)). Let $I^* = \text{Ker}(\pi)$. We now give a multiplicative basis for S. Let \mathcal{B}^* be the set

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \cup \begin{pmatrix} 0 & 0 \\ 0 & b \end{pmatrix} \mid b \in \mathcal{B} \} \cup \begin{pmatrix} 0 & e_i^* b \\ 0 & 0 \end{pmatrix} \mid 1 \le i \le n, b \in \mathcal{B} \}.$$

We can define an order $<^*$ on \mathcal{B}^* by $\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} <^* \begin{pmatrix} 0 & 0 \\ 0 & b \end{pmatrix}$ for each b. $\begin{pmatrix} 0 & 0 \\ 0 & b \end{pmatrix} <^* \begin{pmatrix} 0 & 0 \\ 0 & b' \end{pmatrix}$ if b < b'. $\begin{pmatrix} 0 & e_i^*b' \\ 0 & 0 \end{pmatrix}$ for all b, b' and all i. Finally, $\begin{pmatrix} 0 & e_i^*b \\ 0 & 0 \end{pmatrix} <^* \begin{pmatrix} 0 & e_j^*b' \\ 0 & 0 \end{pmatrix}$ if i < j or if i = j, b < b'.

We leave it to the reader to check that $<^*$ is an admissible order on \mathcal{B}^* . Next we give a set of generators for I^* . Assume I is generated by $\{h_i\}_{i \in \mathcal{I}}$ in R. Then I^* is generated by

$$(*) \qquad \{ \begin{pmatrix} 0 & 0 \\ 0 & h_i \end{pmatrix} \mid i \in \mathcal{I} \} \cup \{ \begin{pmatrix} 0 \sum_{j=1}^n e_j^* a_{i,j} \\ 0 & 0 \end{pmatrix} \mid 1 \le i \le m \}.$$

Note that we view the elements $a_{i,j} \in \text{Span}(\text{NonTip}(I))$ in this setting as elements of R. Again the reader may check that (*) is a generating set for I^* . Under reasonable circumstances, we can find a Gröbner basis for I^* with respect to >*. For example, if R is a path algebra, then S is also a path algebra and hence we can algorithmically find Gröbner bases.

Applying Theorem 2.1, we get $S = I^* \oplus \text{Span}(\text{NonTip}(I^*))$. Remember that $S/I^* = \begin{pmatrix} K & M \\ 0 & R/I \end{pmatrix}$. In particular, M can be identified with $\begin{pmatrix} 0 & M \\ 0 & 0 \end{pmatrix}$. It follows that the elements of $\text{NonTip}(I^*)$ of the form $\begin{pmatrix} 0 & c \\ 0 & 0 \end{pmatrix}$ form a K-basis of M. Noting that if $x \in I$ then $\begin{pmatrix} 0 & 0 \\ 0 & x \end{pmatrix} \in I^*$ and hence $\begin{pmatrix} 0 & e_i^* x \\ 0 & 0 \end{pmatrix} \in I^*$ we conclude that the c's occurring in the K-basis of M are all of the form $\sum_{i=1}^n e_i^* c_i$ with $c_i \in \text{Span}(\text{NonTip}(I))$.

Finally, assuming that a Gröbner basis of uniform elements of I^* can be computed, we note that the normal form of an element of M, $\begin{pmatrix} 0 & c \\ 0 & 0 \end{pmatrix}$, when multiplied by an element of R/I, $\begin{pmatrix} 0 & 0 \\ 0 & b \end{pmatrix}$, yields an element of the form $\begin{pmatrix} 0 & c' \\ 0 & 0 \end{pmatrix}$ with $c' = \sum_{i=1}^{n} e_i^* c_i^*$ with $c_i^* \in \text{Span}(\text{NonTip}(I))$.

We identify R/I with $\operatorname{Span}(\operatorname{NonTip}(I))$ (as usual) and M with the set of elements $c = \sum_{i=1}^{n} e_i^* c_i \in G_0$ where $c_i \in \operatorname{Span}(\operatorname{NonTip}(I))$ with $\begin{pmatrix} 0 & c \\ 0 & 0 \end{pmatrix} \in$ $\operatorname{Span}(\operatorname{NonTip}(I^*))$. Note that such c can in fact be viewed as in F_0 and with this identification, we get a vector space splitting $\sigma : M \to F_0$ of $f_0 : F_0 \to M$ by $\sigma(m) = N(\begin{pmatrix} 0 & g \\ 0 & 0 \end{pmatrix})$ where $f_0(g) = m$.

Stepping back from the above details, what we have is the following, assuming we can compute a Gröbner basis of uniform elements of I^* algorithmically:

- 1. One can compute a K-basis of M inside F_0 algorithmically; namely, the Gröbner basis of I^* with respect to $>^*$.
- 2. One can algorithmically compute the action of R/I on this basis; namely, the computation of the normal form of the product.

In practice, we can use a Gröbner basis of I with respect to > in that the elements $\begin{pmatrix} 0 & 0 \\ 0 & g \end{pmatrix}$ are part of a Gröbner basis of I^* . Once the Gröbner basis for I^* is computed, operations by elements of R/I on M can be performed using the division algorithm.

2.5.2 Projective Resolutions In this section, we use the results of the previous section and Gröbner bases to construct projective resolutions of modules. We need to require more properties of R.

Our goal is to show how to algorithmically construct a projective R/I-resolution of a right R/I-module M. This construction will use both the construction of Gröbner bases, overlap relations and the division algorithm.

In this section we assume

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- 1. R has a multiplicative basis \mathcal{B} and an admissible order >.
- 2. *R* is a finitely generated *K*-algebra and \mathcal{B} is generated by $X = \{b_1, \ldots, b_n\}$.
- 3. Every element of \mathcal{B} is a unique product of elements of X; that is, if $c_1 \cdots c_r = d_1 \cdots d_s$ with $c_i, d_i \in \mathcal{B}$ then r = s and $c_i = d_i$.
- 4. I is an ideal with Gröbner basis \mathcal{G} .

As in the previous section, given a finitely presented right R/I-module M with presentation $f_1: F_1 \to F_0$, we form the ring $\bar{S} = \begin{pmatrix} K & M \\ 0 & R/I \end{pmatrix}$. Viewing \bar{S} as a quotient of

$$S = \begin{pmatrix} K & F_0 \\ 0 & R \end{pmatrix}$$

by I^* , we find a Gröbner basis \mathcal{G}^* of I^* with respect to $>^*$. Recall that $\mathcal{G}^* = \{\begin{pmatrix} 0 & 0 \\ 0 & g \end{pmatrix} | g \in \mathcal{G}\} \cup \{\begin{pmatrix} 0 & h_i^* \\ 0 & 0 \end{pmatrix}\}_{i \in \mathcal{I}}$ where \mathcal{G} is a Gröbner basis for I with respect to > and h_i^* are of the form $\sum_{i=1}^n e_i^* c_i$ with $c_i \in \text{Span}(\text{NonTip}(I))$. We assume each $\sum_{i=1}^n e_i^* c_i$ and each $g \in \mathcal{G}$ are are uniform elements of R. Furthermore, assume F_0 has basis e_1, \ldots, e_n

Note that if $h_i^* = \sum_j e_j^* c_{i,j}$, we define F_1' to be the free R/I-module with basis $\{d_i' \mid i \in \mathcal{I}\}$. Define

$$f_1': F_1' \to F_0$$

by $f'_1(d'_i) = \sum_j e_j c_{i,j}$. We see that $f'_1 : F'_1 \to F_0$ is also a free presentation of M.

Let f'_1 be represented by the matrix $(a_{i,j})$ as in the last section. We also let $f_0: F_0 \to M$ be the canonical surjection. Note that the modules (0 R/I)and (K M) are projective \bar{S} -modules. This follows from $(0 R/I) \simeq \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \bar{S}$ and $(K M) \simeq \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \bar{S}$.

Let U be the right \overline{S} -module $(K \ 0)$ (with the action of \overline{S} given by right matrix multiplication). We have a projective presentation of U given as follows.

$$\left(\begin{smallmatrix} 0 & F_0 \end{smallmatrix}\right) \stackrel{\left(\begin{smallmatrix} 0 & f_0 \end{smallmatrix}\right)}{\longrightarrow} \left(\begin{smallmatrix} K & M \end{smallmatrix}\right) \to U \to 0.$$

Note that the image of $(0 f_0)$ is (0 M). We may continue the resolution using f'_1 as follows:

$$\left(\begin{array}{c} 0 \ F_1' \end{array} \right) \stackrel{\left(\begin{array}{c} 0 \ f_1' \end{array} \right)}{\longrightarrow} \left(\begin{array}{c} 0 \ F_0 \end{array} \right) \stackrel{\left(\begin{array}{c} 0 \ f_0 \end{array} \right)}{\longrightarrow} \left(\begin{array}{c} K \ M \end{array} \right) \rightarrow U \rightarrow 0.$$
 (**)

We show how to algorithmically find F_2 , a free R/I-module and a map $f_2: F_2 \to F'_1$ so that

$$F_2 \xrightarrow{f_2} F_1' \xrightarrow{f_1'} F_0 \xrightarrow{f_0} M \to 0$$

is an exact sequence. Once this is done, we can continue the projective resolution by replacing $F'_1 \stackrel{f'_1}{\to} F_0$ with $F_2 \stackrel{f_2}{\to} F'_1$ and repeat the construction. This will be done by continuing the sequence (**) over \bar{S} and finding the next map in the resolution

$$\left(\begin{array}{cc} 0 & F_2 \end{array}\right) \stackrel{\left(\begin{array}{c} 0 & f_2 \end{array}\right)}{\longrightarrow} \left(\begin{array}{cc} 0 & F_1' \end{array}\right).$$

For this we first describe F_2 . Let \mathcal{O} be the set of all overlap relation of form $o(\alpha, \beta, p^*, q^*)$ where α is of the form $\begin{pmatrix} 0 & h_i^* \\ 0 & 0 \end{pmatrix}$ and β is of the form $\begin{pmatrix} 0 & 0 \\ 0 & g \end{pmatrix}$.

Let $o(\alpha, \beta, p^*, q^*)$ be an overlap relation in \mathcal{O} . Let $\alpha = \begin{pmatrix} 0 & h \\ 0 & 0 \end{pmatrix}$ and $\beta = \begin{pmatrix} 0 & 0 \\ 0 & g \end{pmatrix}$. It follows that from the structure of \mathcal{B}^* that $p^* = \begin{pmatrix} 0 & p \\ 0 & 0 \end{pmatrix}$ for some $p \in \mathcal{B}$ and that $q^* = \begin{pmatrix} 0 & 0 \\ 0 & g \end{pmatrix}$ for some $q \in \mathcal{B}$. Furthermore, we have

$$o(lpha,eta,p^*,q^*) = (1/\mathrm{Tip}(lpha))lpha \cdot q^* - (1/\mathrm{Tip}(eta))p^*eta = \ igg(egin{array}{c} 0 & (1/\mathrm{Tip}(h))hq - (1/\mathrm{Tip}(g))pg \ 0 \end{pmatrix} .$$

Since \mathcal{G}^* is a Gröbner basis, the remainder of $o(\alpha, \beta, p^*, q^*)$ by division by \mathcal{G}^* is 0. Hence

$$\begin{aligned} o(\alpha, \beta, p^*, q^*) &= \\ \sum_i \sum_j \begin{pmatrix} 0 & h_i \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & q_{i,j} \end{pmatrix} + \sum_{g \in \mathcal{G}} \sum_j \begin{pmatrix} 0 & u_{i,j} \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & g \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & v_{i,j} \end{pmatrix} \end{aligned}$$

where $q_{i,j}, u_{i,j}, v_{i,j} \in \mathcal{B}$.

We can now describe F_2 and $f_2: F_2 \to F'_1$. F_2 is the free R/I-module with basis \hat{e}_i , where *i* is indexed by the overlap set \mathcal{O} . To describe $f_2: F_2 \to F'_1$ we need only define $f_2(\hat{e}_i)$ for $i \in \mathcal{O}$. If $i = o(\alpha, \beta, p^*, q^*)$, then, keeping the notation of the previous paragraph,

$$f_2(\hat{e}_i) = \sum_{j \in \mathcal{I}} d'_j q_{i,j}.$$

The above description of the construction of projective resolutions is a variation on the results found in [17] which use results from [5]. The proof of the next result can be obtained by vary the proof found in [17] of Theorem 4.1 to our setup. **Theorem 2.4.** Keeping the above notations,

$$F_2 \xrightarrow{f_2} F_1' \xrightarrow{f_1'} F_0 \xrightarrow{f_1} M \to 0$$

is an exact sequence.

2.5.3 Finite Dimensional Algebras and Modules In this section, we show how to find projective presentations from "standard" representations of modules and vice versa. R will denote a finitely generated K-algebra with multiplicative basis \mathcal{B} , and generators b_1, \ldots, b_n and admissible order >. Let I be an ideal such that R/I is finite dimensional. Let M be a finite dimensional right R/I-module. As usual, we identify R/I with Span(NonTip(I)) and let $\text{Span}(\text{NonTip}(I)) = \overline{\mathcal{B}}$. Without loss of generality, we may assume that the generators b_1, \ldots, b_n of \mathcal{B} are in $\overline{\mathcal{B}}$. (If some $b_i \notin \overline{\mathcal{B}}$ then b_i acts as 0 on M and the following can be modified accordingly.)

Let $\{m_j\}_{j \in \mathcal{J}}$ be a K-basis for M. Well-order the m_j 's arbitrarily. Let $D = |\mathcal{J}|$. Viewing M as a D-dimensional vector space, each generator b can be represented by a $D \times D$ -matrix $(c_{i,j}^b)$ with entries in K, where

$$m_i \cdot b = \sum_{j=1}^D m_j c_{i,j}^b.$$

The $D \times D$ matrices $(c_{i,j}^{b_l})$, l = 1, ..., n might be called the "standard" way of representing M.

First we show how to obtain the $(c_{i,j}^{b_l})$ from an R/I-free presentation of $M, f_1 : F_1 \to F_0$. Suppose F_0 has R/I-basis $\{e_1, \ldots, e_r\}$ and F_1 has R/I-basis $\{d_1, \ldots, d_s\}$. Then, viewing f_1 as a K-linear map M is the cokernel of f_1 . Note that the K-bases of F_0 and F_1 are just $\{e_i b \mid 1 \leq i \leq r \text{ and } b \in \overline{B}\}$ and $\{d_i b \mid 1 \leq i \leq s \text{ and } b \in \overline{B}\}$ respectively. Hence knowledge of f_1 as a R/I-homomorphism easily yields viewing f_1 as a K-linear map. Furthermore, for each generator b_l we have a commutative diagram:

$$\begin{array}{cccc} F_1 & \stackrel{f_1}{\to} & F_0 \\ \downarrow b_l & & \downarrow b_l \\ F_1 & \stackrel{f_1}{\to} & F_0 \end{array}$$

where the vertical morphisms are right multiplication by b_l . It follows that we get a map $g_{b_l}: M \to M$ representing multiplication by b_l . g_{b_l} yields the matrix $(c_{i,j}^{b_l})$.

To go the other way, assume that we have a K-basis $\{m_j\}_{j\in\mathcal{J}}$ of M and the $D \times D$ K-matrices $(c_{i,j}^{b_l})$ representing multiplication by the n generators $b_l, 1 \leq l \leq n$. We show how to construct an R/I-projective presentation $f_1: F_1 \to F_0$ of M. First let F_0 be the free R/I-module with basis $\{e_j \mid j \in \mathcal{J}\}$ and F_1 be the free R/I-module with basis $\{d_{x,j} \mid x \in \overline{\mathcal{B}}, j \in \mathcal{J}\}$. Define $f_0: F_0 \to M$ by $f_0(e_j) = m_j$. Next, define $f_1: F_1 \to F_0$ as follows. Given $j \in \mathcal{J}$ and $x \in X$,

since x is a product of b_l 's, using the matrices $(c_{u,v}^{b_l})$, we may calculate $m_j \cdot x = \sum_{i \in \mathcal{J}} m_i \alpha_{i,j}^x$ where $\alpha_{i,j}^x \in K$. We now define $f_1(d_{x,j}) = e_j x - \sum_{i \in \mathcal{J}} e_i \alpha_{i,j}^x$.

Proposition 2.12. Keeping the notations above,

$$F_0 \xrightarrow{f_1} F_1 \xrightarrow{f_0} M \to 0$$

is exact.

Proof. We have seen that $f_0(F_0) = M$. Hence the sequence is exact at M. It is immediate that $f_0f_1(d_{x,j}) = 0$ for each $x \in \overline{\mathcal{B}}$ and $j \in \mathcal{J}$. Thus $\operatorname{Im}(f_1) \subset \operatorname{Ker}(f_0)$. It remains to show that $\operatorname{Ker}(f_0) \subset \operatorname{Im}(f_1)$.

Consider $\{z = \sum_{j \in \mathcal{J}} e_j w_j \in \operatorname{Ker}(f_0) | z \notin \operatorname{Im}(f_1)\}$. We wish to show this set is empty. Let $z = \sum_{j \in \mathcal{J}} e_j w_j$ be in this set. Suppose $w_j = \sum_{x \in \overline{\mathcal{B}}} \beta_{j,x} x$. Consider $z' = z + f_1(\sum_{j,x} d_{j,x}\beta_{j,x})$. Since $z \notin \operatorname{Im}(f_1), z' \notin \operatorname{Im}(f_1)$. But, from the definition of f_1 , we see $f_1(\sum_{j,x} d_{j,x}\beta_{j,x}) = \sum_{j \in \mathcal{J}} e_j \gamma_j - z$ for some $\gamma_j \in K$. Thus $z' = \sum_{j \in \mathcal{J}} e_j \gamma_j$. But f(z') = 0 and hence $\sum_{j \in \mathcal{J}} m_j \gamma_j = 0$. It follows that $\gamma_j = 0$ for all $j \in \mathcal{J}$. This is a contradiction and we have shown no such z can exist.

2.6 Applications of Gröbner Bases

Throughout this section, we will tacitly assume that $R = K\Gamma$, \mathcal{B} is the set of paths, and > is an admissible order. Hence Gröbner bases exist for ideals in R.

2.6.1 Topics not covered in Detail

- 1. Computation of $\operatorname{Hom}_{R/I}(M, N)$ for two right R/I-modules can easily be performed. More precisely, Gröbner bases yield K-bases of right modules which behave "nicely" with respect to the ring structure of R/I. That is, viewing R/I as Span(NonTip(I)) with "rewrite formulas" given by the Gröbner basis, the module bases constructed in the last lecture come with how the basis of R/I acts on basis elements of the module because that is precisely the information contained in the Gröbner basis of I^* in $\binom{K \ G_0}{0 \ R}$ of the last lecture. Using these bases, the computation of $\operatorname{Hom}_{R/I}(M, N)$ as a vector space is straightforward.
- 2. The study of submodules and quotient modules can be easily handled using the bases of last lecture.
- 3. Computation of $M \otimes_{R/I} N$ for a left R/I-module M and a right R/Imodule N can be performed. For this one must be able to find bases of both left and right modules (which Gröbner bases can do). One method of attacking the computation $M \otimes_{R/I} N$ is to find $M \otimes_K N$ using the K-bases provided by the Gröbner bases of the last lecture and then put in the relations over R/I and find quotient.

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- 4. Using the above two parts, it should be apparent that one can computationally study module theory of noncommutative rings that are quotients of rings that have a theory of Gröbner bases. This class includes algebras that are quotients of path algebras and therefore includes quotients of free algebras.
- 5. One can study the growth of algebras. By Theorem 2.1, R/I can be identified with Span(NonTip(I)) and NonTip(I) is a K-basis of R/I. If R is the free algebra and \mathcal{B} is the set of monomials or if R is a path algebra and \mathcal{B} is the set of paths then the growth of the number of nontips of a given degree or length measures the growth of the algebra. More precisely, if $H_n = \dim_K(\text{Span}(\{x \in \text{NonTip}(I) \mid \text{length}(x) = n\})$, then the Hilbert series of R/I is

$$H(R/I) = \sum_{n=o}^{\infty} H_n z^n.$$

Here z is a variable. Questions like the rationality of H(R/I) are being studied [2, 3, 4, 30].

One of the important tools to study H(R/I) and the numbers H_n is the Ufnarovski graph. Space doesn't permit me to say much about it, but it is a powerful tool to study the nontips of I. We refer to [32].

- 6. Noncommutative Gröbner bases in free algebras have been applied to study H^{∞} control problems in the work of Helton, Stankus and Varvik.
- 7. As described in the last section, Gröbner bases allows the construction of projective resolutions of a right module M. Then, applying $\operatorname{Hom}_{R/I}(-, N)$ to the resolution and taking cohomology of the resulting complex, one can obtain the ext-groups $\operatorname{Ext}_{R/I}^n(M, N)$. For a left R/I-module N, applying $-\otimes_{R/I} N$ to the resolution and taking cohomology of the resulting complex, one can obtain the tor-groups $\operatorname{Tor}_n^{R/I}(M, N)$. Hence, one can study homological questions computationally.

One important invariant of an algebra is the *Poincaré series of a module*. If M is a right R/I-module and S is a simple right R/I-module (usually K), let $P_n = \dim_K \operatorname{Ext}^n_{R/I}(M, S)$. The Poincaré series is

$$P_S(M) = \sum_{n=0}^{\infty} P_n z^n.$$

Again, questions like rationality and connections with the Hilbert series are of interest. [4, 30]

8. Michael Bardzell [8] has used the theory of Gröbner basis to study the Hochschild cohomology of a monomial K-algebra.

2.6.2 Algebras and their associated Monomial Algebras One of the interesting connections one obtains from the theory of Gröbner bases is the associated monomial algebra to an algebra. As we have done, let $I_{MON} = \text{Span}(\text{Tip}(I))$. Then I_{MON} is a monomial ideal. We say R/I_{MON} is the associated monomial algebra to R/I. Note that R/I_{MON} is dependent on the choice of basis \mathcal{B} of R and on the choice of admissible order >.

If R/I is a finite dimensional K-algebra, there is an important invariant of R/I called the *Cartan matrix*. Let S_1, \ldots, S_n be a full set of nonisomorphic simple R/I-modules and let P_1, \ldots, P_n be indecomposable projective R/I-modules such that there are surjections $P_i \to S_i$. Then we define the Cartan matrix of R/I to be the $n \times n$ matrix $(C_{i,j})$ where $C_{i,j}$ is the number of times S_i occurs as a composition factor of P_j .

Since R is a path algebra, it is sometimes natural to assume that I is contained in the ideal $\langle \Gamma_1 \rangle^2$ where Γ_1 is the set of arrows in Γ . In this case there are certain simple modules of R/I we distinguish, called the *vertex simple modules*. A vertex simple module is a module of the form $S_i = v_i R / \langle \Gamma_1 \rangle$ where v_i is a vertex. It is easy to see that if there are n vertices in Γ , S_1, \ldots, S_n are right simple R/I-modules which are 1-dimensional over K.

If M is a right R-module and

$$\cdots P_n \xrightarrow{f_n} P_{n-1} \to \cdots \to P_1 \xrightarrow{f_1} P_0 \xrightarrow{f_0} M \to 0$$

is a projective resolution of M, then the projective dimension of M, $pd_{R/I}(M)$, is the smallest n such that $Im(f_n)$ is a projective module. If no such n exists, we say that M has infinite projective dimension and write $pd_{R/I}(M) = \infty$. The global dimension of R/I, gl. dim(M), is N if

 $N = \sup\{ pd_{R/I}(M) \mid M \text{ is an } R/I\text{-module} \}$

if such an N exists. Otherwise, gl. $\dim(M) = \infty$.

We can state a result relating R/I and R/I_{MON} .

Theorem 2.5. [23] Let $R = K\Gamma$ be a path algebra with K-basis \mathcal{B} the finite directed path in Γ and with admissible order >. Let I be an ideal in R and I_{MON} be the associated monomial ideal. Then

- 1. $H(R/I) = H(R/I_{MON})$.
- 2. If R/I is finite dimensional, then

$$\dim_K(R/I) = \dim_K(R/I_{MON}) = |NonTip(I)|.$$

- 3. Assume that $I \subset <\Gamma_1>^2$.
 - (a) The construction of a projective resolution given in [17] for each vertex simple module S_i is minimal for R/I_{MON} .
 - (b) For each $i = 1, \ldots, n$, $\operatorname{pd}_{R/I}(S_i) \leq \operatorname{pd}_{R/I_{MON}}(S_i)$.

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- (c) If R/I is finite dimensional, then gl. dim $(R/I) \leq$ gl. dim (R/I_{MON}) .
- (d) If R/I is finite dimensional, then the Cartan matrices of R/I and R/I_{MON} are equal.
- (e) If S is a vertex simple module viewed as a module over both R/I and R/I_{MON} , and if $\rightarrow P_1 \rightarrow P_0 \rightarrow S \rightarrow 0$ and $\rightarrow P'_1 \rightarrow P'_0 \rightarrow S \rightarrow 0$ are minimal projective resolutions over R/I and R/I_{MON} respectively, then $\dim_K(P_n) \leq \dim_K(P'_n)$ for all $n \geq 0$.

2.6.3 Koszul Algebras Let $R = K\Gamma$ and I be an ideal in R. We say that I is a quadratic ideal if there is a set of generators $\{f_1, \ldots, f_n\}$ of I such that each f_1 is a K-linear combination of paths in Γ of exactly length 2. In this case, we say that $\{f_1, \ldots, f_n\}$ is a set of quadratic generators and that R/I is a quadratic algebra.

The path algebra R has a natural positive **Z**-grading given by letting vertices be homogeneous of degree 0, and paths be homogeneous of degree equal to the length. We will view R as a graded ring in this way. Note that an element $x = \sum_{i=1}^{m} \alpha_i p_i$ with $\alpha_i \in K$ and $p_i \in \mathcal{B}$ is homogeneous if each p_i occurring in x has the same length. We call this the length grading of R. Let I be a graded ideal in R; that is, an ideal generated by homogeneous elements. Then R/I is a **Z**-graded algebra by the induced grading.

For the remainder of this section, we assume that I is graded ideal in R and let $R/I = S_0 \oplus S_1 \oplus S_2 \oplus \cdots$ as a graded ring. We let $\operatorname{Gr}(S)$ denote the category of \mathbb{Z} -graded S-modules and degree 0 S-module maps. That is, the objects of $\operatorname{Gr}(S)$ are

$$M = \cdots \oplus M_{-2} \oplus M_{-1} \oplus M_0 \oplus M_1 \oplus M_2 \oplus \cdots$$

where each i, j, M_i is a right S_0 -module and if $s_j \in S_j$ and $m_i \in M_i$ then $m_i s_j \in M_{i+j}$ such that, forgetting the graded structure, M is a right S-module. A degree 0 map $f: M \to N$ between graded S-modules is an S-module map such that if $m_i \in M_i$ then $f(m_i) \in N_i$.

We say a graded S-module is generated in degree n if $M_j = 0$ for j < n and for all $i \ge 0$, the multiplication maps $M_n \otimes_{S_0} S_i \to M_{n_i}$ are surjective. We say that an S-module X is gradable if there is a graded S-module M, such that X is isomorphic to M when one forgets the grading on M and views M as an S-module. We have the following result whose proof is standard.

Proposition 2.13. Let I be a graded ideal in a path algebra $K\Gamma$ where $K\Gamma$ has the length grading. Assume that $I \subset <\Gamma_1 >$. Let $K\Gamma/I = S_0 \oplus S_1 \oplus S_1 \oplus \cdots$ be the graded quotient ring. Then

- 1. S_0 is isomorphic to $K\Gamma/\langle \Gamma_1 \rangle$ and hence is semisimple.
- 2. Each S_i is finite dimensional over K.
- 3. The vertex simple modules are gradeable modules.
- 4. $S_1 \oplus S_2 \oplus \cdots$ is the graded Jacobson radical of S.

- 5. The category of finitely generated graded S-modules has projective covers.
- 6. Graded S-modules have graded projective resolutions.
- 7. A graded S-projective resolution of a graded S-module M forgets to a Sprojective resolution of X where X is the S-module M when we forget the graded structure.

We say R/I is Koszul algebra if S_0 , viewed as a graded S-module generated in degree 0 has a graded projective resolution

$$\cdots P_n \to P_{n-1} \to \cdots \to P_1 \to P_0 \to M \to 0$$

such that for each n, P_n is generated in degree n. Koszul algebras are an important class of algebras that naturally occur in algebraic geometry, topology, and the theory of quantum groups [10, 12, 14, 29, 31, 33]. Proofs of many of the basic results about Koszul algebras can be found in [21, 22, 10]. My goal here is to demonstrate that Gröbner bases can be used in studying such a class of Koszul algebras.

It is well-known that if R/I is a Koszul algebra then I must be a quadratic ideal [21, 10]. At this time, there is no classification known of which quadratic ideals I have the property that $K\Gamma/I$ is a Koszul algebra at this time, in terms of the generators of the quadratic ideals. We do have the following result though.

Theorem 2.6. [20] Let I be a quadratic ideal in a path algebra $K\Gamma$. Let > be an admissible order on the paths such that I has a quadratic Gröbner basis. Then $K\Gamma/I$ is a Koszul algebra.

The proof of the result, although too technical for these lectures, involves an investigation of the projective resolution of the vertex simple modules given by the construction discussed in the last lecture. Analysis of the construction shows that if the algebra is graded and the module is graded, then the constructed resolution is, in fact, a graded resolution of the module. It is also shown that if the generators of I are quadratic, then the construction of a graded projective resolution of a vertex simple module has the desired degree properties.

The next result follows from the above theorem and also from [23].

Corollary 2.3. Let I be a monomial ideal generated by some paths of length 2. Then $K\Gamma/I$ is a Koszul algebra.

Proof. If J is a monomial ideal in a path algebra, since overlap relations are in fact 0, we see by the Termination Theorem that any generating set of monomials for J is a Gröbner basis under any admissible order. Since I can be generated by paths of length 2, it follows that I has a quadratic Gröbner basis.

We end with another application which is more fully described in [19]. Let $R = K < x_1, \ldots, x_n >$ be the free associative algebra in n noncommuting variables. Let > be the degree-lexicographic order with $x_1 < x_2 < \cdots < x_n$. For $1 \le i < j \le n$, let

$$q_{i,j} = x_j x_i - c_{i,j} x_i x_j + r_{i,j}$$

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where $r_{i,j}$ is a quadratic polynomial, each of whose terms is less than $x_i x_j$. Thus $\operatorname{Tip}(q_{i,j}) = x_j x_i$ and $\operatorname{Tip}(q_{i,j} - \operatorname{Tip}(q_{i,j}) = x_i x_i$. Let I be the ideal generated by $\{q_{i,j}\}$.

Note that I is a quadratic ideal. Consider R/I. We denote the image of x_i in R/I by \bar{x}_i . We say R/I has a *Poincaré-Birkhoff-Witt basis* or PBW basis if $\{\bar{x}_1^{a_1}\bar{x}_2^{a_2}\ldots\bar{x}_n^{a_n}\}$ where a_i are nonnegative integers, is a K-basis of R/I.

The set $\{q_{i,j}\}_{1 \le i,j \le n}$ can be viewed as rewriting rules in the sense that if a monomial $\bar{m} = \bar{x}_{i_1} \bar{x}_{i_2} \dots \bar{x}_{i_s}$ has the property that for some $j, i_j > i_{j+1}$ then in $R/I, \bar{x}_{i_j} \bar{x}_{i_{j+1}} = c_{i_j,i_{j+1}} \bar{x}_{i_j} - \bar{r}_{i_j,i_{j+1}}$,

$$\bar{m} = \bar{x}_{i_1} \bar{x}_{i_2} \dots \bar{x}_{i_{j-1}} \bar{x}_{i_{j+1}} \bar{x}_{i_j} \bar{x}_{i_{j+2}} \dots \bar{x}_{i_s} - \bar{x}_{i_1} \bar{x}_{i_2} \dots \bar{x}_{i_{j-1}} \bar{r}_{i_j} \bar{x}_{i_{j+2}} \dots \bar{x}_{i_s}$$

Thus, by replacing $\bar{x}_j \bar{x}_i$ as above if j > i and by the fact that > is a wellorder, we see that $\{\bar{x}_1^{a_1} \bar{x}_2^{a_2} \dots \bar{x}_n^{a_n}\}$ generate R/I as a K vector space. It is natural to ask if R/I has a PBW basis. The next result answers the question.

Proposition 2.14. Keeping the above notations, R/I has a PBW basis if and only if $\{q_{i,j}\}$ is a Gröbner basis of I. Thus, if R/I has a PBW basis, R/I is a Koszul algebra.

Proof. Suppose that R/I has a PBW basis. Then $\{\bar{x}_1^{a_1}\bar{x}_2^{a_2}\ldots\bar{x}_n^{a_n}\}$ is a K-basis for R/I. To show that $\{q_{i,j}\}$ is a Gröbner basis for I it suffices to show that overlap relations have remainder 0 when divided by $\{q_{i,j}\}$.

Consider division by $\{q_{i,j}\}$. Any monomial m occuring in a remainder cannot have $x_j x_i$ with j > i in it since if so, $\operatorname{Tip}(q_{i,j})$ would divide m. Thus, the remainder of division by $\{q_{i,j}\}$ of any overlap relation is in the span of $\{x_1^{a_1} \dots x_n^{a_n}\}$. By the PBW basis assumption, no element of this span is in I other than 0.

If $\{q_{i,j}\}$ is a Gröbner basis for I, then

NonTip $(\{q_{i,j}\}) = \{m \in \mathcal{B} \mid \text{Tip}(q_{i,j}) \text{ does not divide } m\}.$

But it immediate that NonTip $(\{q_{i,j}\})$ is $\{x_1^{a_1} \dots x_n^{a_n} | a_i \ge 0\}$. But NonTip(I) is a K-basis of R/I under the usual identification. Thus, R/I has a PBW basis.

The last result of the proposition is a consequence of Theorem 2.6. $\hfill \Box$

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